



Education

# NAND Flash Solid State Storage Performance and Capability -- an In-depth Look

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## **NAND Flash Solid State Storage Performance and Capability**

"This presentation provides an in-depth examination of the fundamental theoretical performance, capabilities, and limitations of NAND Flash-based Solid State Storage (SSS). The tutorial will explore the raw performance capabilities of NAND Flash, and limitations to performance imposed by mitigation of reliability issues, interfaces, protocols, and technology types. Best practices for system integration of SSS will be discussed. Performance achievements will be reviewed for various products and applications. "

# Moore's continues to beat Newton's Law

- ◆ **Mechanical Drives have hit their limits**
  - ◆ Platter stability degrades at higher speeds
  - ◆ Short-stroking reduces capacity for seek time
  - ◆ Capacity is limited by smaller form factors
- ▶ **Solid State Storage continues to evolve**
  - ◆ Greatest bit density (bits per cubic volume)
  - ◆ Random IOPS are 250 times greater
  - ▶ MLC increases capacity and lowers costs
  - ◆ Advanced error correction improves reliability
  - ◆ Performance and Capacity are intertwined



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**There can be no data integrity trade-off  
for performance**

# Media Reliability / Availability

## ➤ The GOOD

- ◆ No moving parts
- ◆ Post infant mortality (catastrophic) device failures are rare
- ◆ Predictable wear out

## ◆ The BAD

- ◆ Relatively high bit error rate, which increases with wear
- ◆ Higher density and MLC increases bit error rate
- ◆ Program and Read Disturbs

## ◆ The UGLY

- ◆ Partial Page Programming
- ◆ Data retention is poor at high temperature and wear
- ◆ Infant mortality is high (large number of parts...)

# Controller Reliability Management **SNIA**

- Wear leveling & Spare Capacity
- ◆ Read & Program Disturb control
- ◆ Data & Index Protection
  - ECC Correction
  - ◆ Internal RAID
  - ◆ Data Integrity Field (DIF)
- ◆ Management

**Poor Media + Great Controller = Great SSS Solution**



# Data Integrity versus Performance



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# Performance is about ROI

## Lower OpEx

- Less HW Maintenance
- Less SW Maintenance
- ◆ Greater Uptime
- ◆ Less Power/Cooling
- Fewer Diverse Skills

## Lower CapEx

- Fewer CPUs
- Less RAM
- ◆ Less Network Gear
- ◆ Fewer SW Licenses
- Less Space

**HIGHER Productivity**

# Media Performance

## ➤ The GOOD

- ◆ Performance is excellent (wrt HDDs)
- ◆ High performance per power (IOPS/Watt)
- ◆ Low pin count: shared command / data bus → good balance

## ◆ The BAD

- ◆ Not really a random access device
  - Block oriented
  - R/W access speed imbalance
    - ◆ Slow effective write (erase/transfer/program) latency
- ◆ Performance changes with wear

## ◆ The UGLY

- Some controllers do read/erase/modify/write
- ◆ Others use inefficient garbage collection

# Performance Drivers – SSS Design

- Number of NAND Flash Chips (Die)
- Number of Channels (Real / Pipelined)
- ◆ Interconnect
- ◆ Data Protection (internal/external RAID; DIF; ECC...)
- SLC / MLC Flash Type
- ◆ Effective Block Size (LBA; Sector)
- ◆ Write Amplification Efficiency
- Garbage Collection (GC) Efficiency
- ◆ Bandwidth Throttling
- ◆ Buffer Capacity & Mgmt

# Simplified Theoretical Analysis

## ➤ Bandwidth Only (Not IOPS)

- ◆ Large Transfers (Data length = Integer times die count)
- ◆ Infinite Buffer
- ◆ Reads/Writes queued for maximum bandwidth
- ◆ No system latency

## ◆ Read/Write Ratio %'s fixed

- ◆ 100/0, 75/25, 50/50, 25/75, 0/100
- Steady State, 100% Efficient GC (EB erase/EB written = 1)

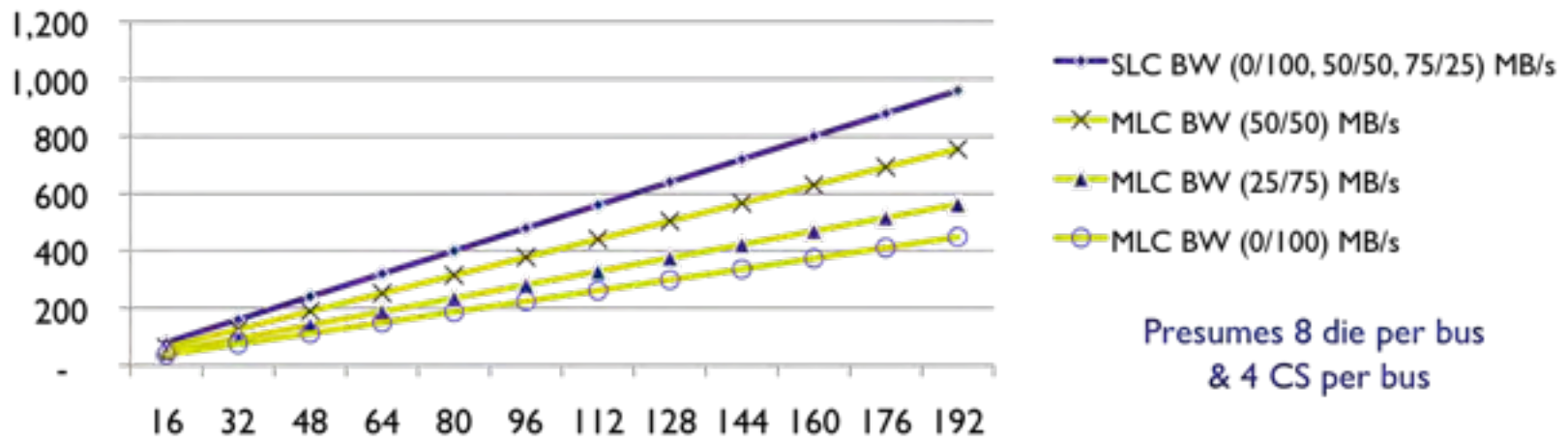
## ◆ Maximum Total BW for SATA-II and PCI-e X4

- ◆ No overhead considered

# Bandwidth Depends on Die Count

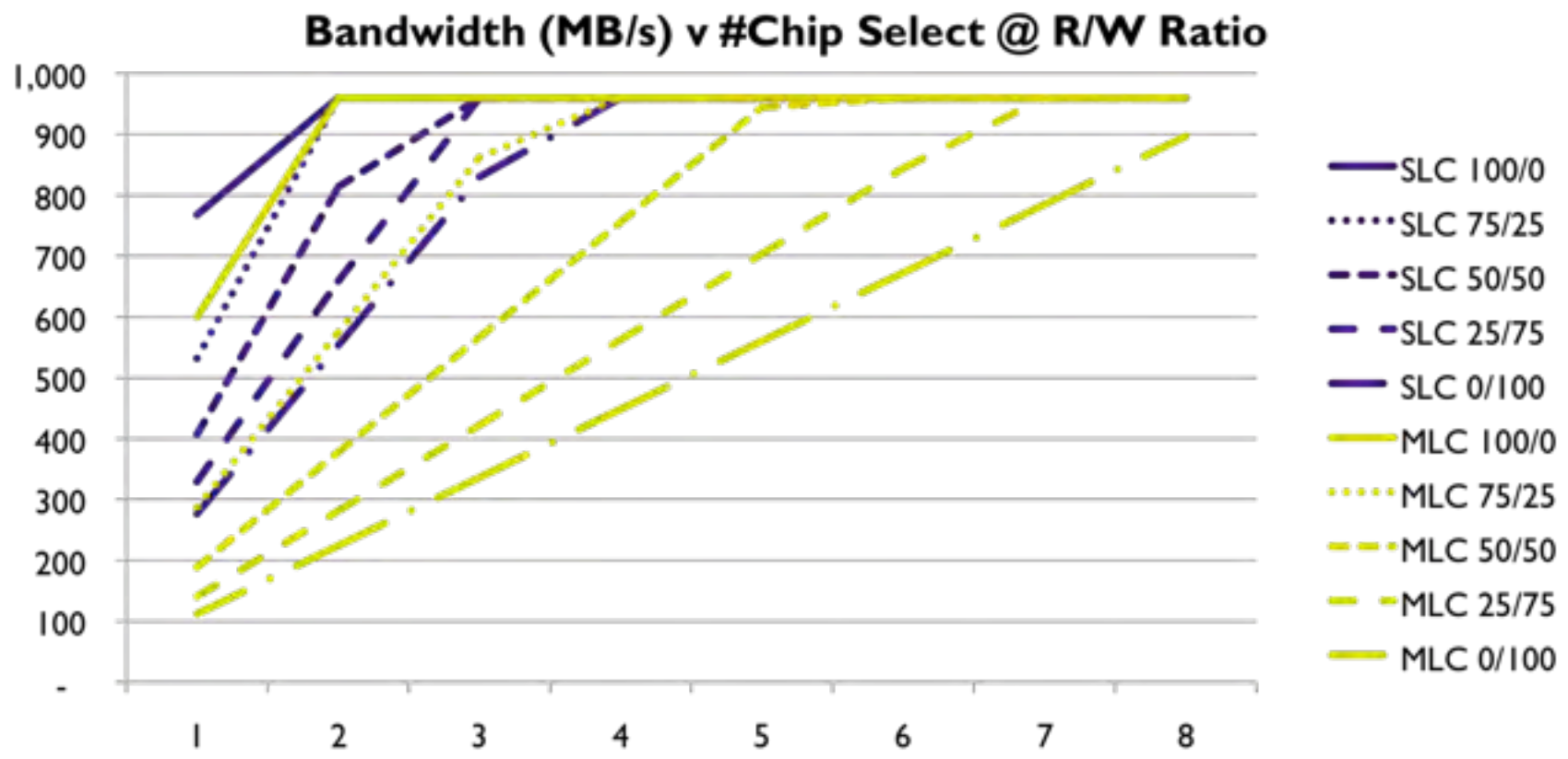
		SLC	MLC
<b>Transfer Rate (MB/s)</b>	tRC & tWC	400	400
<b>Page Program (us)</b>	tProgram	200	600
<b>EB Erase (us)</b>	tErase	3000	10,000
<b>Load Page (us)</b>	tR (tRead)	25	60
<b>Capacity per die</b>		0.5	1.0

**Theoretical BW (MB/s) v Number of Die (SLC, MLC)**



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# Single-Level versus Multi-Level Cells



**Read / write performance imbalance closed with additional banks**  
**Greater R/W imbalance in MLC requires more banks**

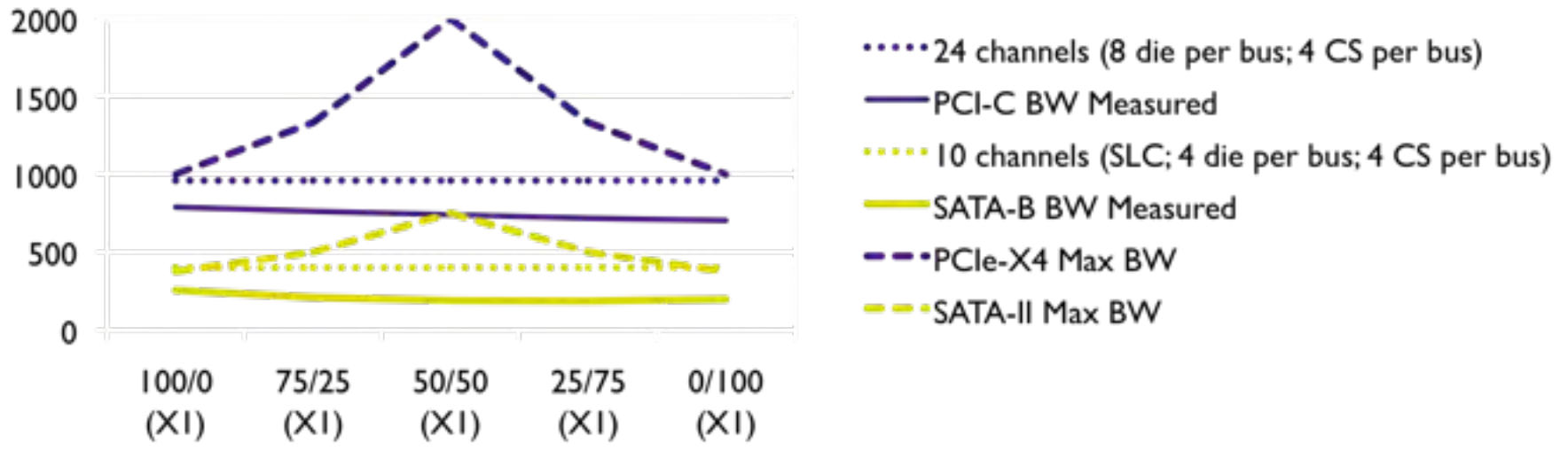
## Features directly affecting performance measurements

	SATA (A)	SATA (B)	PCI (C)
<b>Capacity (GB)</b>	32	32	160
<b>Bus/Link</b>	SATA-II (3 Gb/s)	SATA-II (3 Gb/s)	PCI-E X4 1.1
<b>Memory Type</b>	SLC	SLC	SLC
<b>Adjustable Reserve Capacity</b>	No	No	Yes
<b>SSS Internal RAID</b>	No	No	Yes
<b>-- Running during test</b>	N/A	N/A	Yes
<b>K-IOPS (RMS)</b>	8	27	88
<b>K-IOPS (RMS) / WATT</b>	3	?	7
<b>Bandwidth (RMS, MB/s)</b>	56	208	743
<b>ECC correction</b>	7 bits in 512B	4 bits in ?	11 bits in 240B



# Measured vs Theoretical Bandwidth

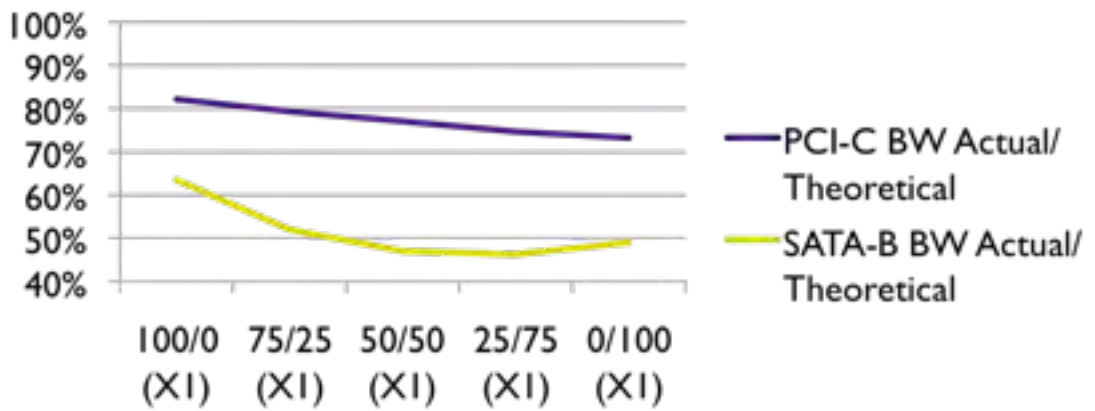
## Measured versus Theoretical Max BW



Note: Theoretical Max BW with 24 channels (4 die per bus, 4 CS per bus) is identical to the PCI-C, 24 channel shown in these charts.

Capacity Multiplier:  
 SATA-B: 1  
 PCI-C: 2

## Measured BW as % of Theoretical Max



## ➤ Read Access

- ◆ Address Chip / EB / Page
- ◆ Load Page into Register
- ◆ Transfer Data From Register 1-byte per cycle

## ➤ Write Access

- ◆ Address Chip / EB
- ◆ Erase EB

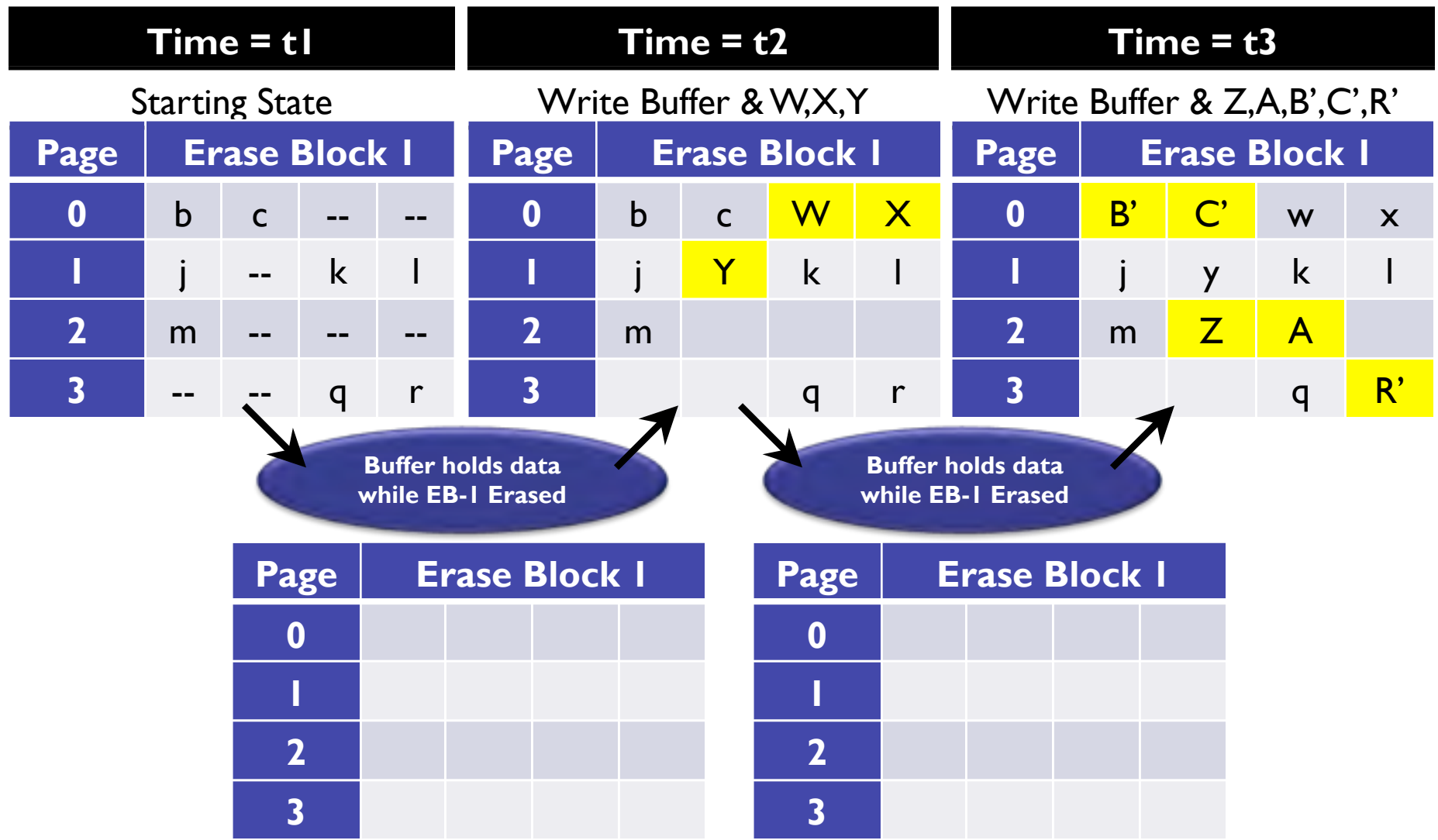
## ◆ ...some time later...

- ◆ Address Chip / EB / Page
- ◆ Transfer Data To Register 1-byte per cycle
- ◆ Program Register to Page

### Typical NAND Flash Die:

- 2000 Erase Blocks (EB)
- 64 Pages per EB
- 4000 Bytes per Page
- 500 MByte Total Capacity

# Example 1: Read/Erased/Modify/Write



# Example 2: Read/Modify/Write

**Time = t1**

Starting State

Page	Erase Block 1			
0	b	c	--	--
1	j	--	k	l
2	m	--	--	--
3	--	--	q	r

**Time = t2**

Data to Buffer (not shown)

Erase EB-1 (not shown)

Write Buffer & W,X,Y to EB-1

Page	Erase Block 2			
0	b	c	W	X
1	j	Y	k	l
2	m			
3			q	r

**Time = t3**

Data to Buffer (not shown)

Erase EB-1 (not shown)

Write Z,A & Replace b,c,r with B',C',R' & Write EB-1

Page	Erase Block 3			
0	B'	C'	w	x
1	j	y	k	l
2	m	Z	A	
3			q	R'

Implicit wear leveling; EB-1 → EB-2 → EB-3

Presumes that destination EB-2 & EB-3 erased prior to transfer of data → higher performance (than previous “Read/Erase/Modify/Write” example)

# Example 3: Garbage Collection

**Time = t1**

Start Garbage Collect EB-1

Page	Erase Block 1			
0	b	c	--	--
1	j	--	k	l
2	m	--	--	--
3	--	--	q	r

Page	Erase Block 2			
0				
1				
2				
3				

**Time = t2**

EB-1 GC'd to EB-2

W,X,Y added

Page	Erase Block 1			
0	b	c	--	--
1	j	--	k	l
2	m	--	--	--
3	--	--	q	r

Page	Erase Block 2			
0	W	b	c	X
1	Y	j	k	l
2	m	q	r	
3				

**Time = t3**

EB-1 erase

b,c,r replaced by B',C',R'

Page	Erase Block 1			
0				
1				
2				
3				

Page	Erase Block 2			
0	w	--	--	x
1	y	j	k	l
2	m	q	--	--
3	B'	C'	R'	--

# GC Performance Impact

## ➤ In this example,

- ◆ COPIED DATA: {b, c, j, k, l, m, q, r} 8 blocks
- ◆ NEW DATA {W, X, Y, B', C', Z, A, R'} 8 blocks

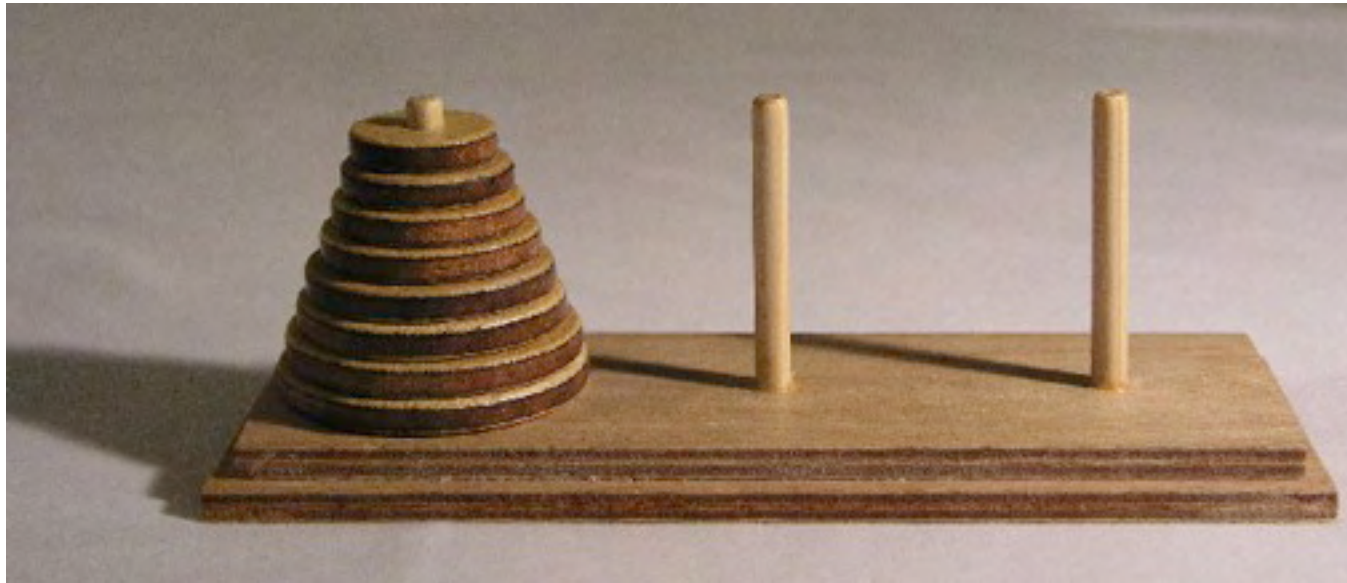
- ◆ 50% (8 of 16) writes are user initiated
- ◆ 50% (8 of 16) writes are internal movement (overhead)

## ◆ Important:

- ◆ 50% of EB-I was “invalid data”
- What if only 10% had been “invalid data?”
- ◆ GC efficiency is dependent upon % of reserve capacity



# Tower of Hanoi



Want to do this in fewer moves?  
Add more pegs!

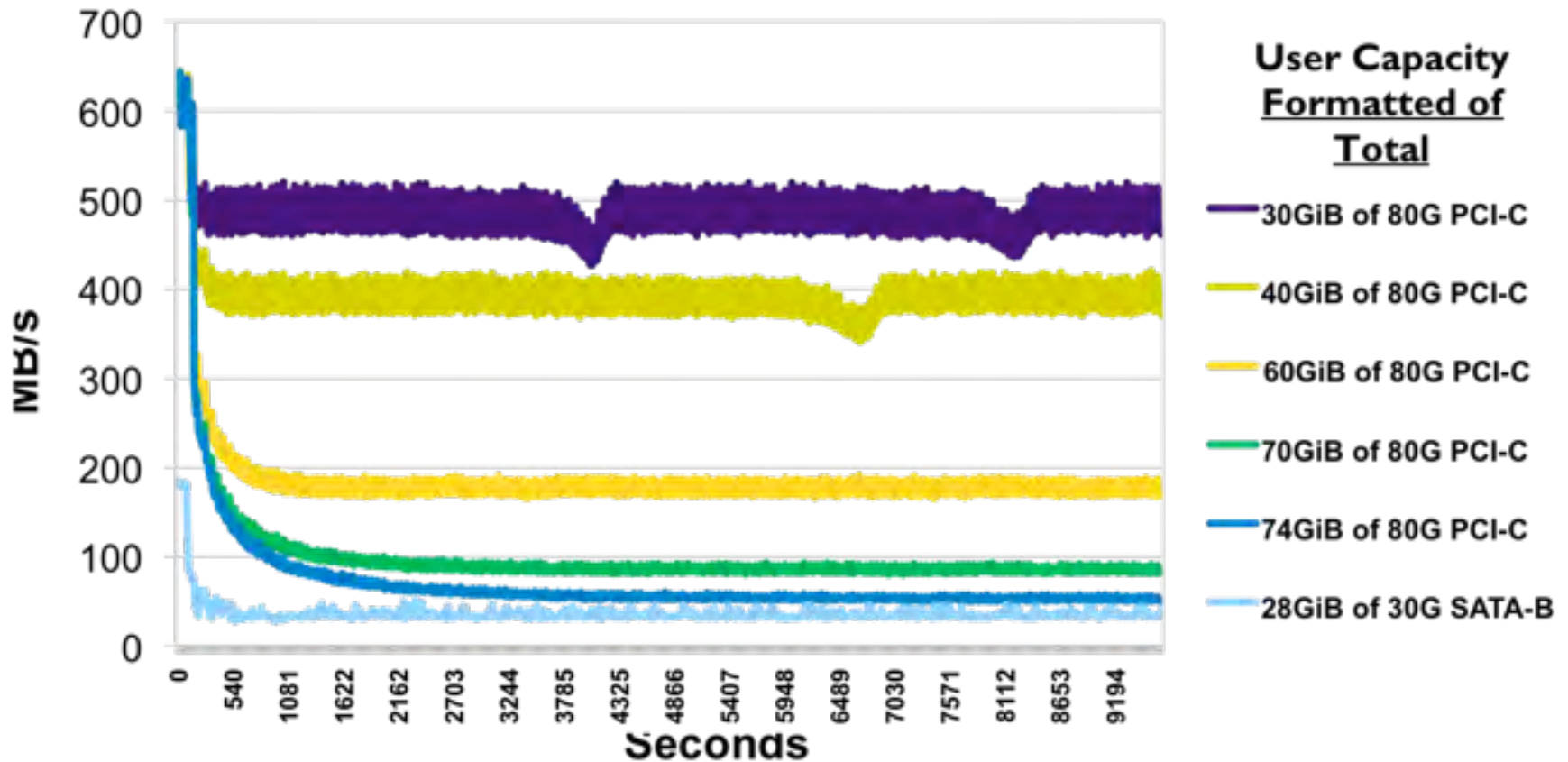
# GC: Pathological Write Conditions

- If a high percentage of total storage capacity utilized  
**AND**
- ◆ A High percentage of data has no correlation-in-time  
**AND**
- ◆ Continuous writing (no recovery time for GC)  
**THEN...**

***Efficiency of GC greatly diminished***

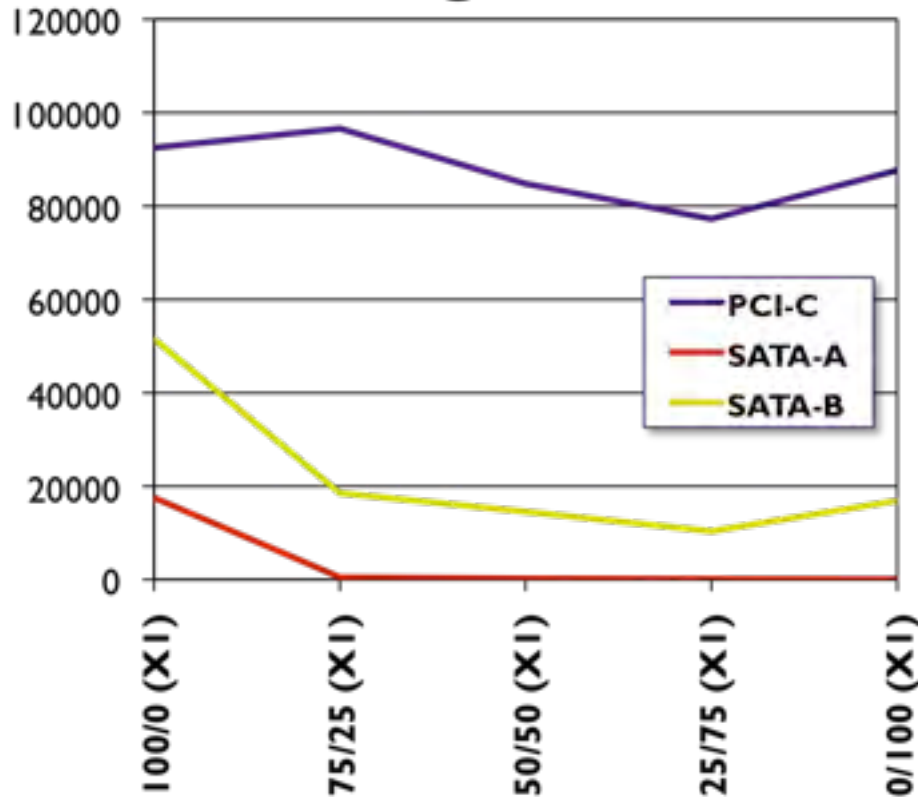


# Pathological Write Condition

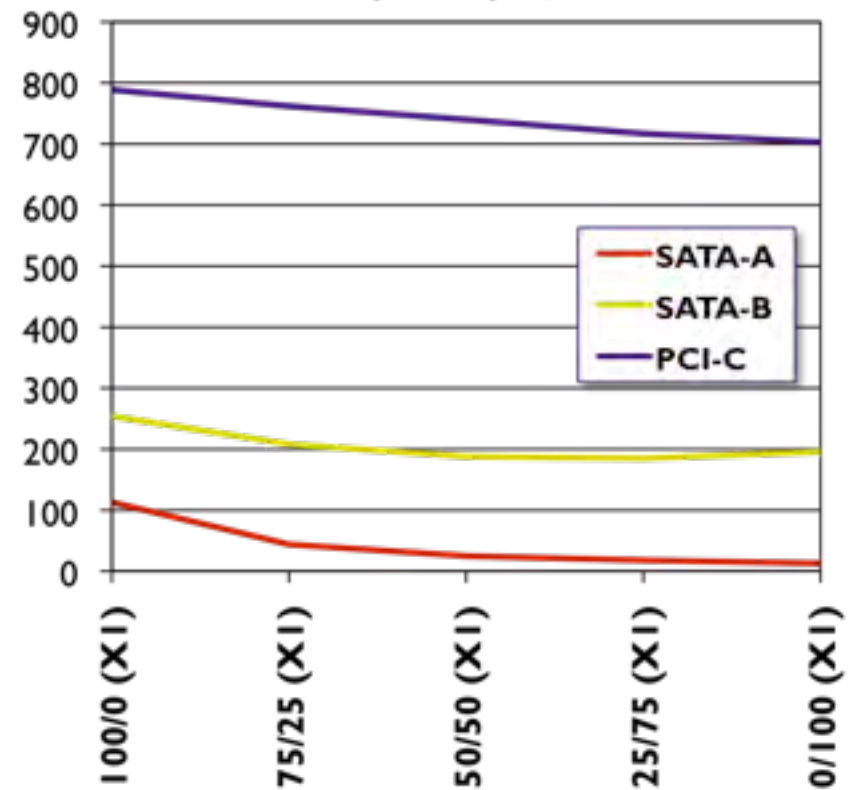


# Performance vs R/W Ratio

**IOPS @ 512 B**

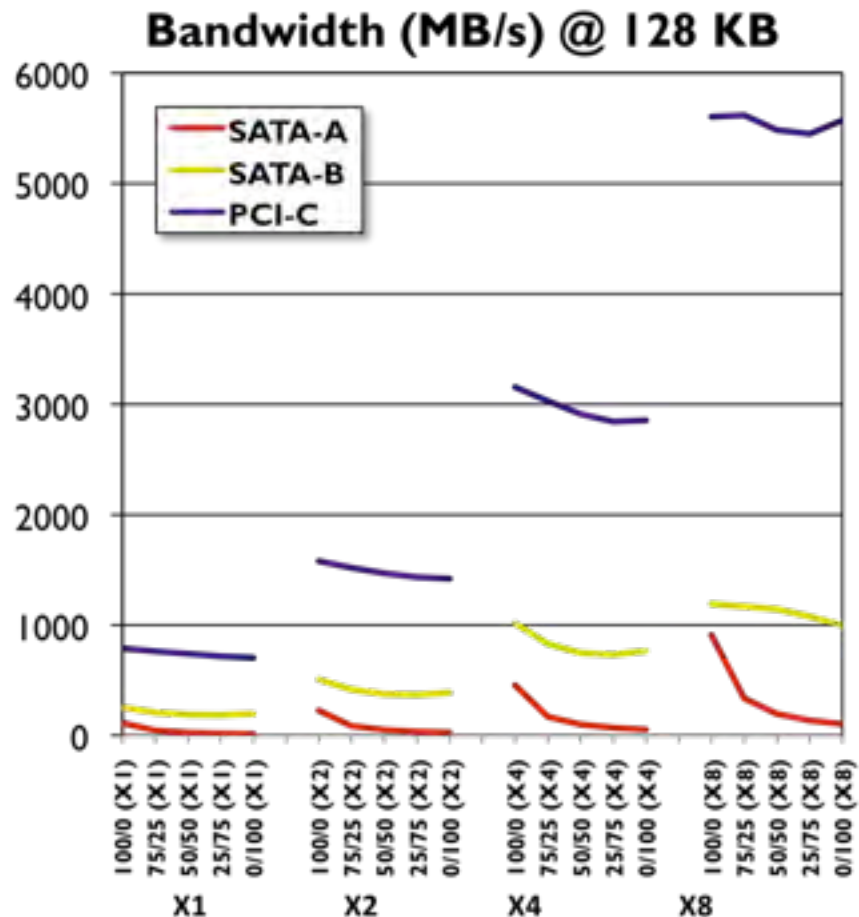
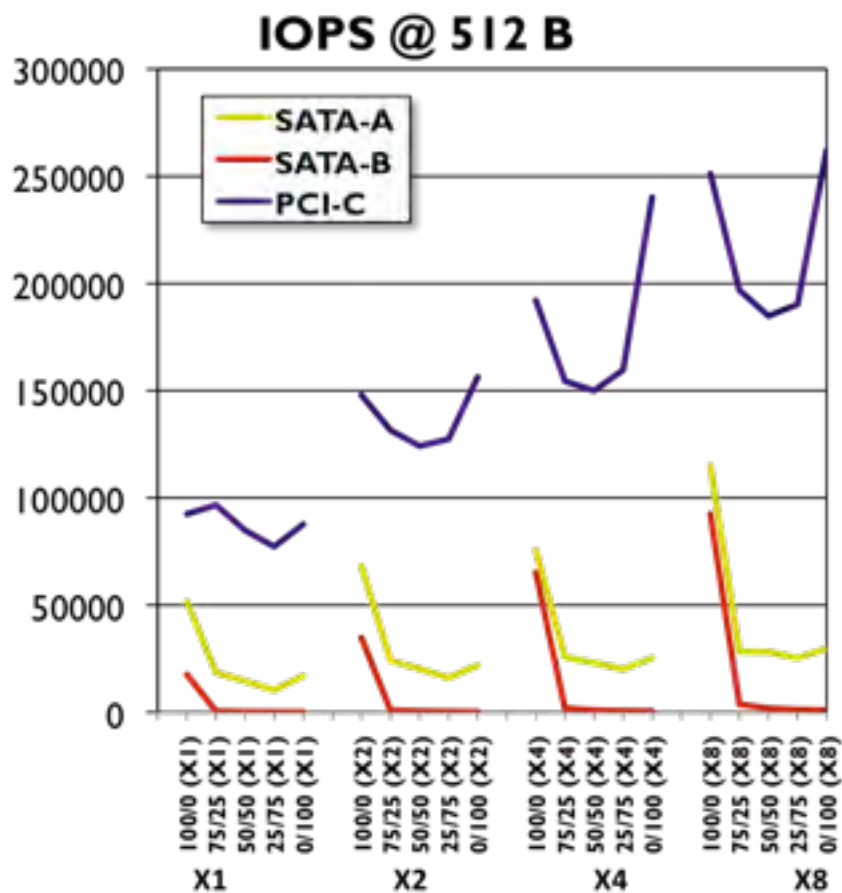


**Bandwidth (MB/s) @ 128 KB**



**Read/Write Collisions → Drop in Mixed Performance**

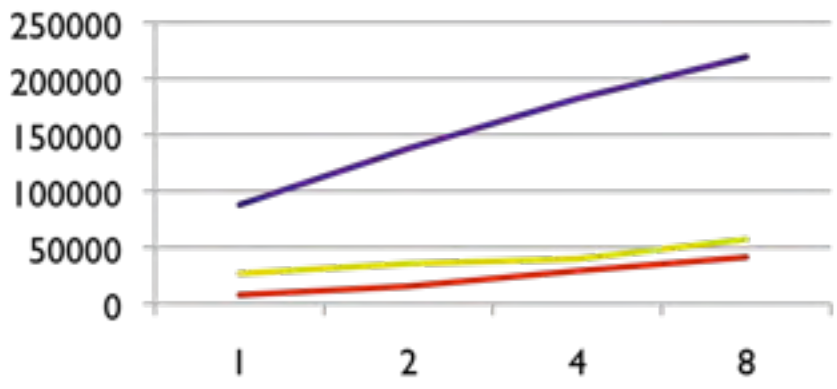
# Scalability versus R/W Ratio



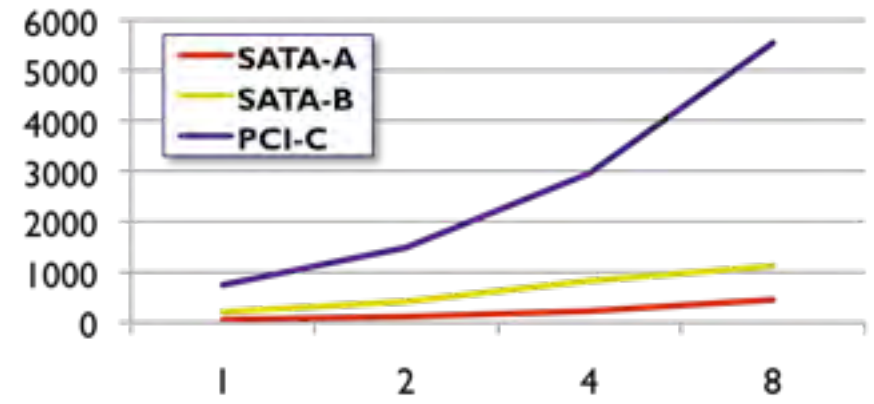
R/W Ratio and Number of Devices in Parallel

# RMS Scalability (# SSS Units)

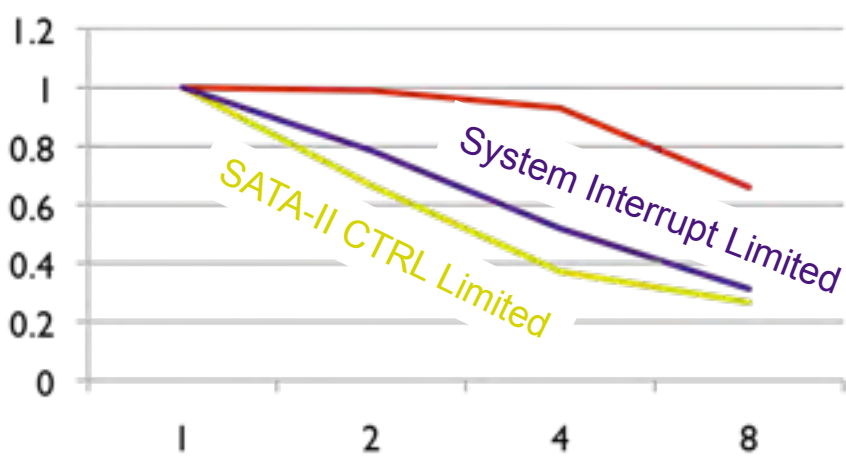
**RMS of IOPS v Scale**



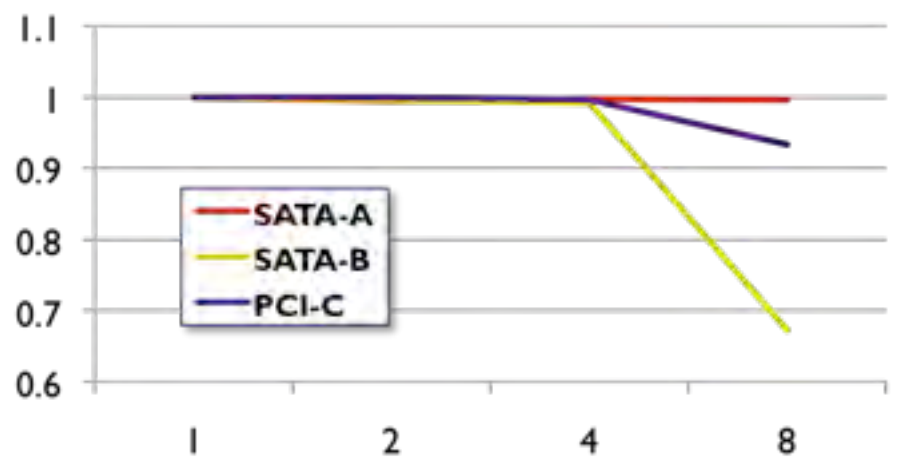
**RMS of Bandwidth v Scale**



**Normalized RMS IOPS v Scale**

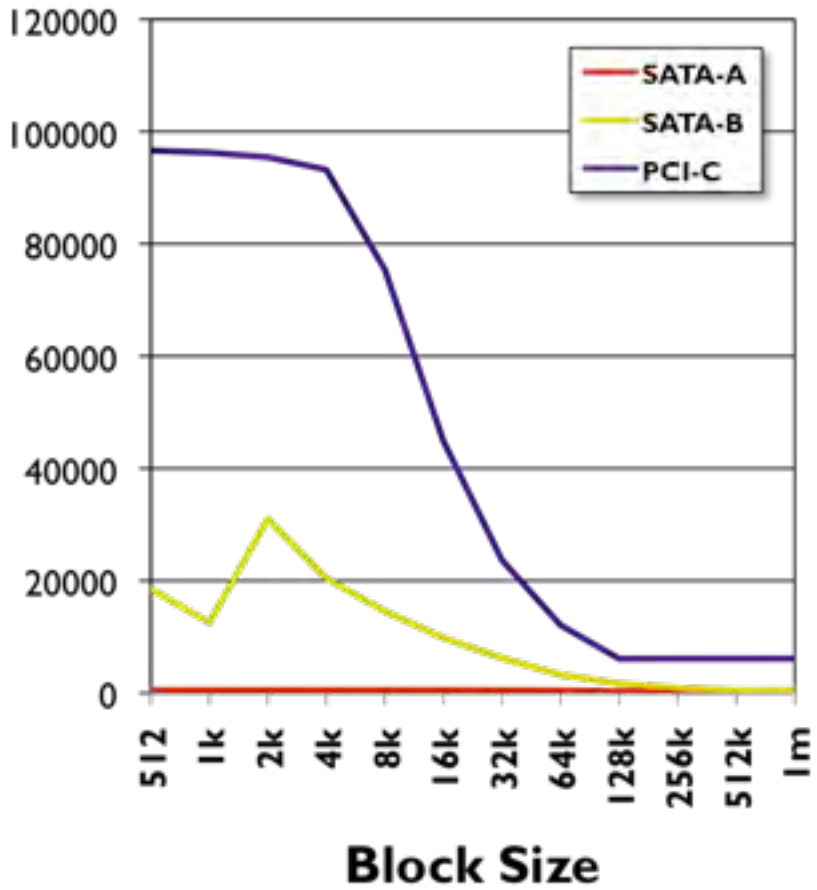


**Normalized RMS Bandwidth v Scale**

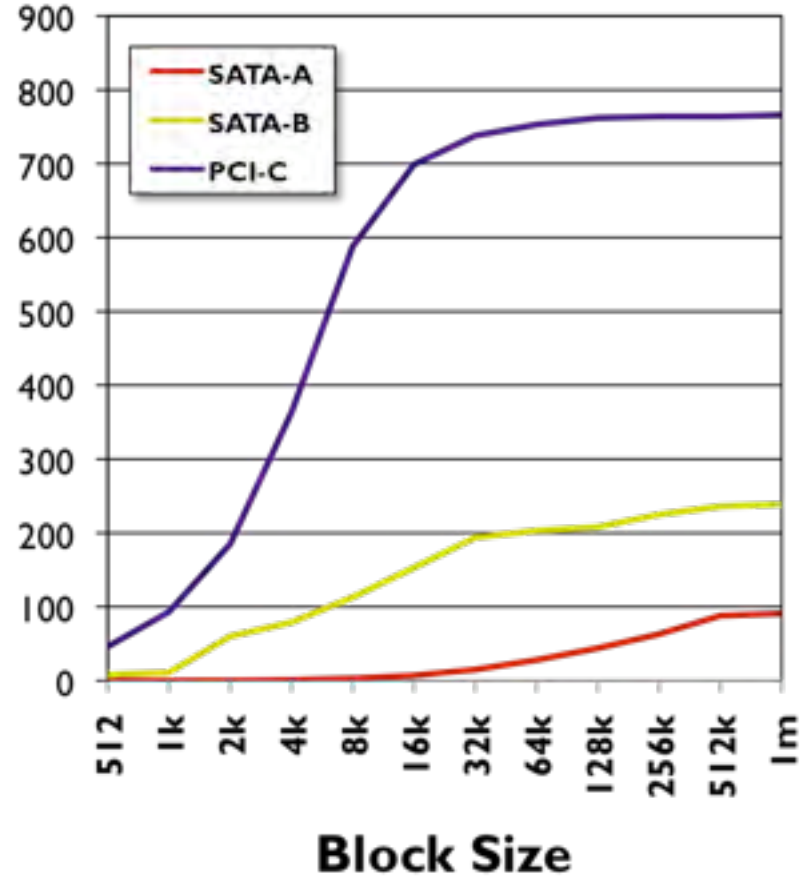


# Performance vs Block Size (75/25) SNIA

**75/25 R/W IOPS**

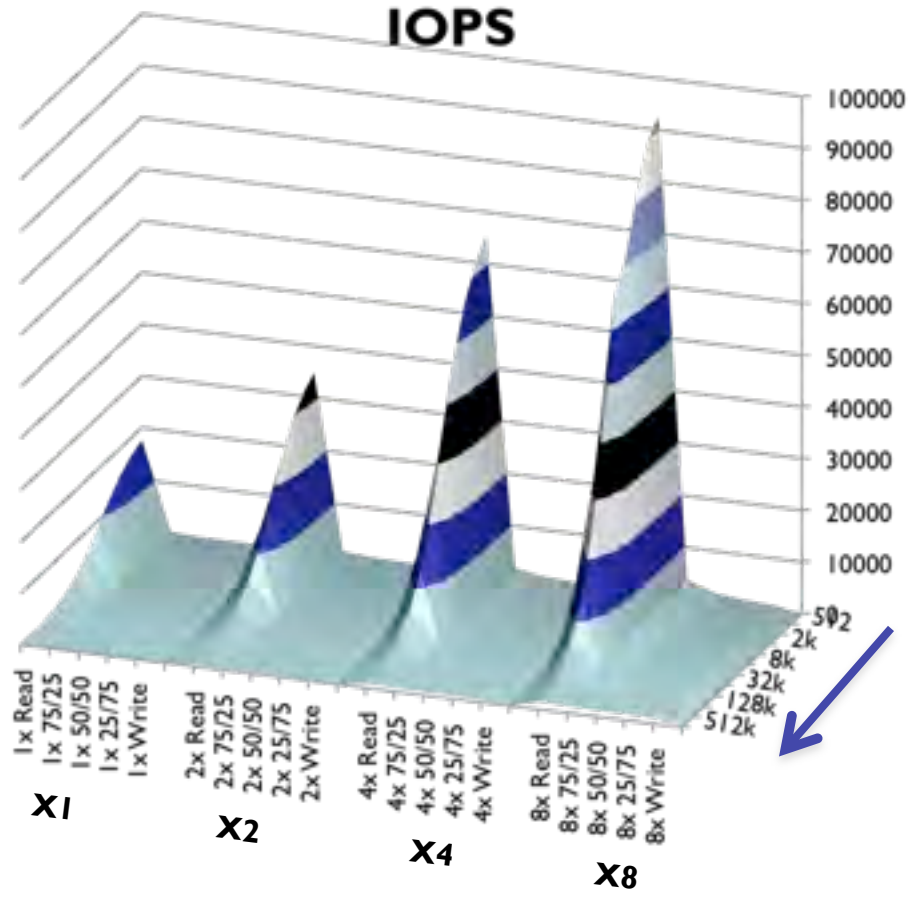


**75/25 R/W Bandwidth (MB/s)**

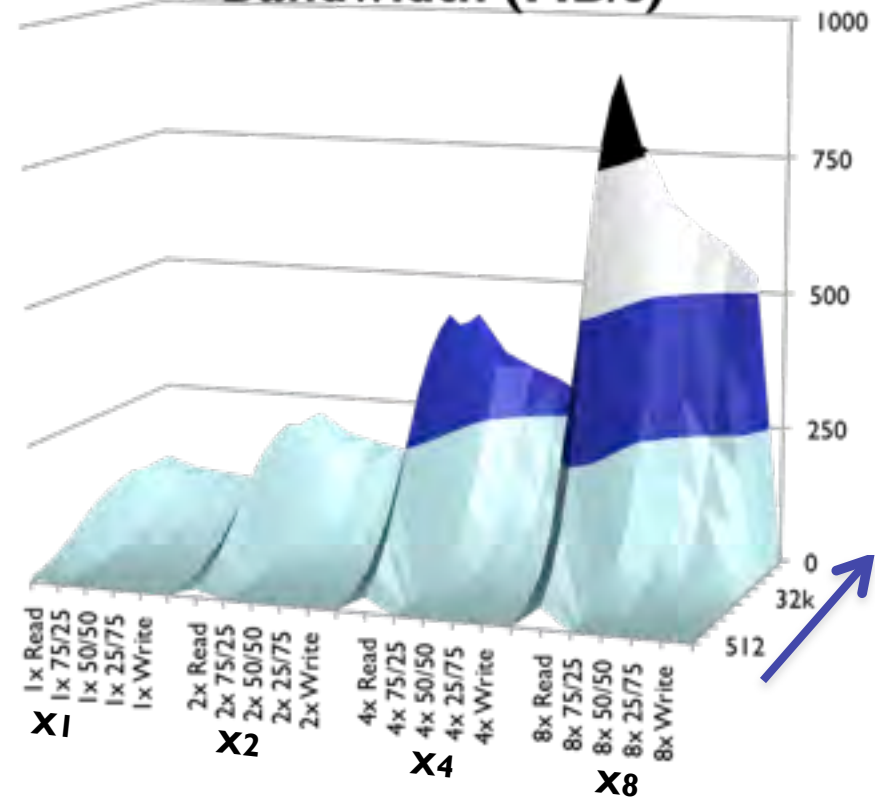


# SATA-A Scalability vs R/W vs Block Size

### SATA-A Scalability IOPS



### SATA-A Scalability Bandwidth (MB/s)

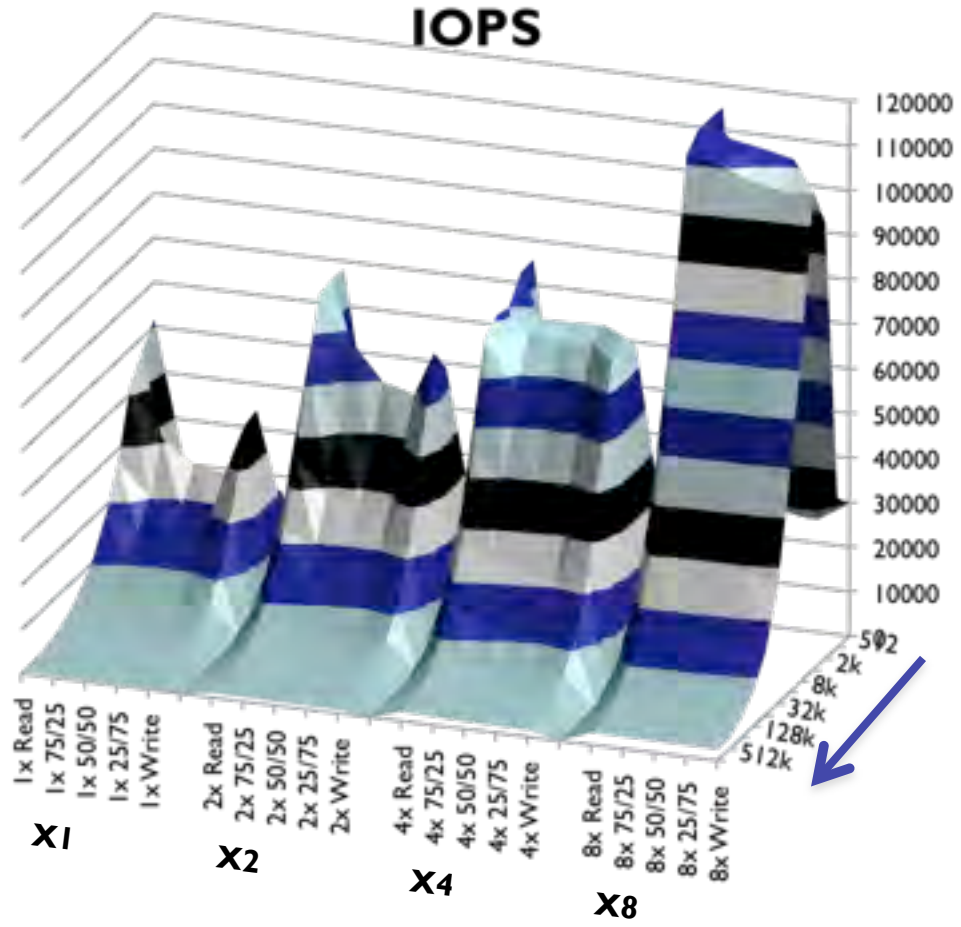


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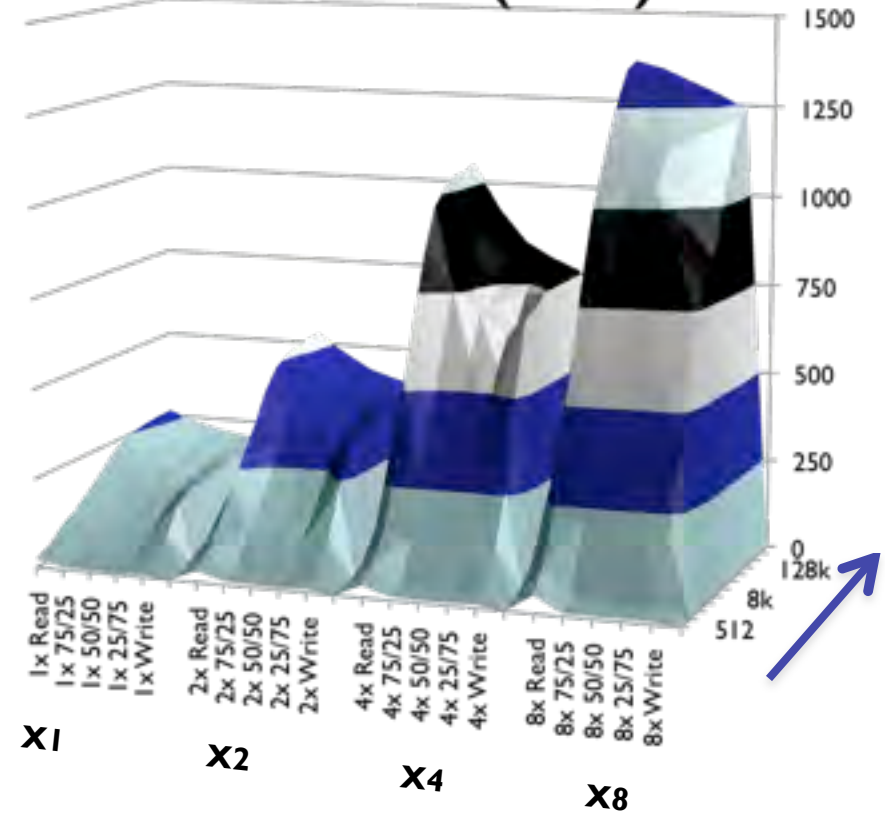


# SATA-B Scalability vs R/W vs Block Size

## SATA-B Scalability IOPS

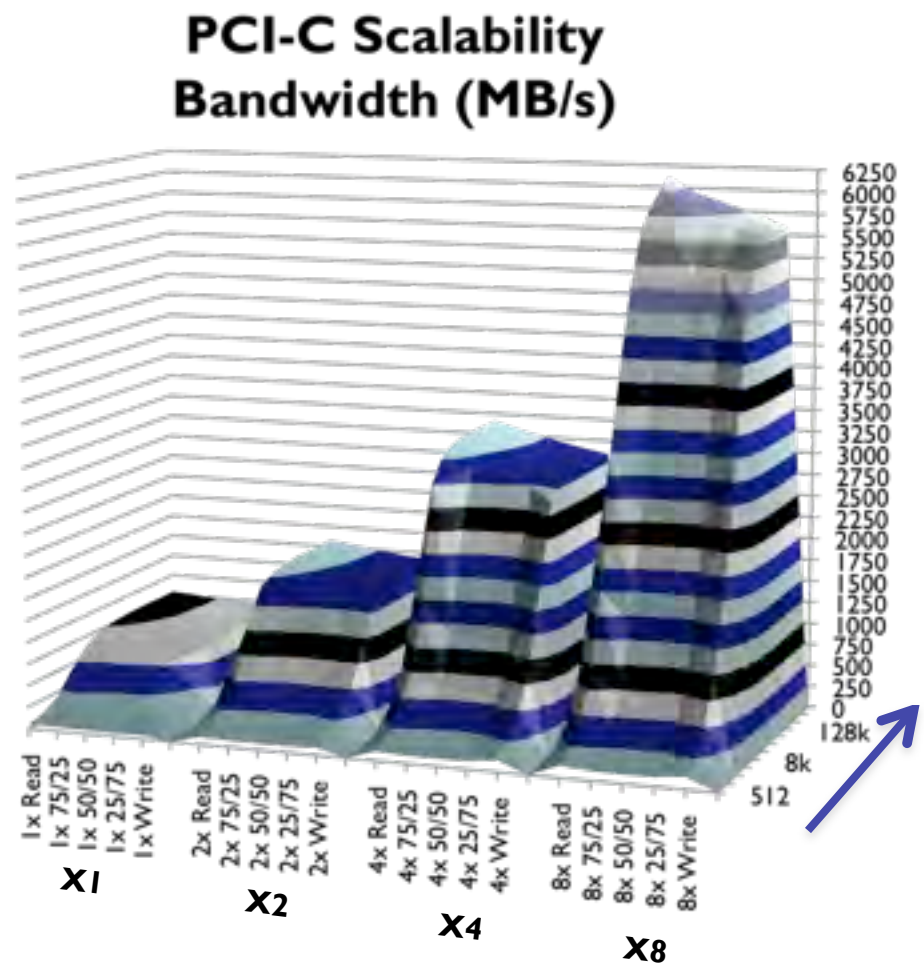
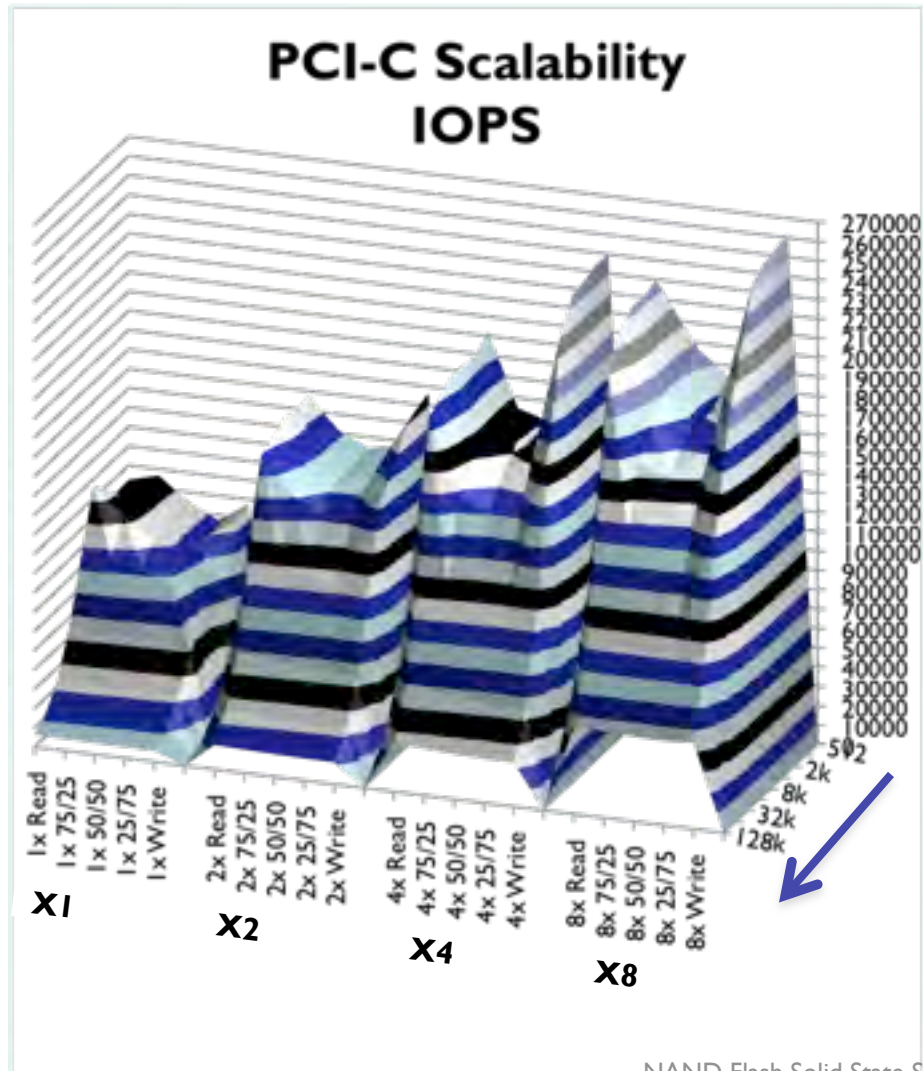


## SATA-B Scalability Bandwidth (MB/s)



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# PCI-C Scalability vs R/W vs Block Size



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# System Level Considerations

- Data / Index Protection (RAID and DIF)
- Scalability
- ◆ Compare system- or data-center-level
  - ◆ Not device
- ◆ Best case: test on real application
  - ◆ Not benchmark
  - ◆ Plan to do tuning to reach top perf. / objectives
  - ◆ Applications may have contra-indicated optimizations
    - > Keeping data in close physical proximity (short stroking)
    - > Caching algorithms

# Questions to Ask : Things to Know

## ☒ Bandwidth / IOPS at

- ☒ **Block size(s) you need**
- ◆ **R/W ratio you use**
- ◆ **Steady State / Burst**
- ☒ Reserve capacity used
- ◆ Data's temporal relationship
- ◆ Scalability
- ☒ RAIDing
- ◆ BOL / EOL

## ☒ Design impacts on data integrity; life; failures & perf.

- ☒ ECC robustness
- ◆ Write amplification / GC efficiency
- ◆ Internal RAID
- ☒ Bandwidth throttling
- ◆ Partial Page Programming

## ◆ Test Conditions

- ◆ Workload
- ◆ Temporal Relationships
- ☒ User capacity / reserve capacity

- ✉ Please send any questions or comments on this presentation to SNIA: : [\*\*tracksolidstate@snia.org\*\*](mailto:tracksolidstate@snia.org)

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